On Applying The λs_e-style of Unification for Simply-Typed Higher Order Unification in the Pure λ-Calculus

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Talk's Plan

1. HOU in explicit substitutions calculi
2. HOU in the λs_e-calculus
   - 2.1. Unification in the λs_e-style of explicit substitutions
   - 2.2. Checking arithmetic constraints (versus shifts and composition in $\lambda\sigma$)
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1. HOU in explicit substitution calculi

HOU $\left\{ \begin{array}{c} \text{Given two simply-typed lambda terms } a \text{ and } b \\ \text{find a } \textit{substitution } \theta \text{ such that} \\ \theta(a) =_{\beta\eta} \theta(b) \end{array} \right.$

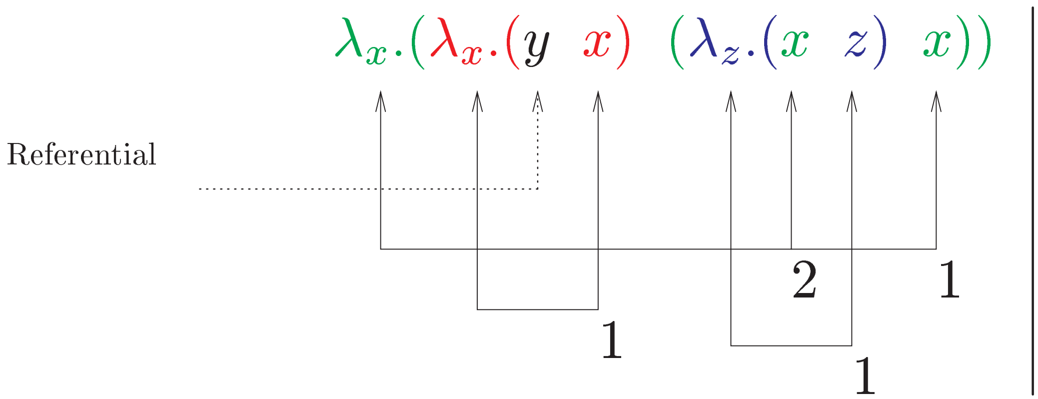
- HOU essential for generalizations of the Robinson's first-order resolution principle.
- HOU applied in
  - Automated (Higher order) reasoning
  - Higher order proof assistants
  - Higher order logic programming

Why *making substitutions explicit* is adequate for reasoning about HOU?

- Substitution is the key operation for HOU.
- *Implicitness* of substitution is the "Achilles heel" of the λ-calculus:
  - β-reduction is given via informal/implicit variable renaming
Implicit substitution does not provide any formal mechanism for analysing essential computational properties

  such as
  - – time and
  - – space complexity

- Terms in de Bruijn notation, $\Lambda_{dB}(\mathcal{X})$: $a ::= \mathbb{N} \mid \mathcal{X} \mid (a \;\; a) \mid \lambda.a$, where $\mathcal{X}$ meta-variables and $\mathbb{N}$ set of de Bruijn indices.



$\rightarrow_\beta \;\; \lambda_x.(y \;\; (\lambda_z.(x \;\; z) \;\; x))$

$\rightarrow_\beta \;\; \lambda_x.(y \;\; (x \;\; x))$

For instance, for the referential $x, y, z, \ldots$: $\lambda.(\lambda.(4 \;\; 1) \;\; (\lambda.(2 \;\; 1) \;\; 1))$

β-reduction: $\lambda.(\lambda.(4 \;\; 1) \;\; (\lambda.(2 \;\; 1) \;\; 1)) \rightarrow_\beta \lambda.(3 \;\; (\lambda.(2 \;\; 1) \;\; 1)) \rightarrow_\beta \lambda.(3 \;\; (1 \;\; 1))$

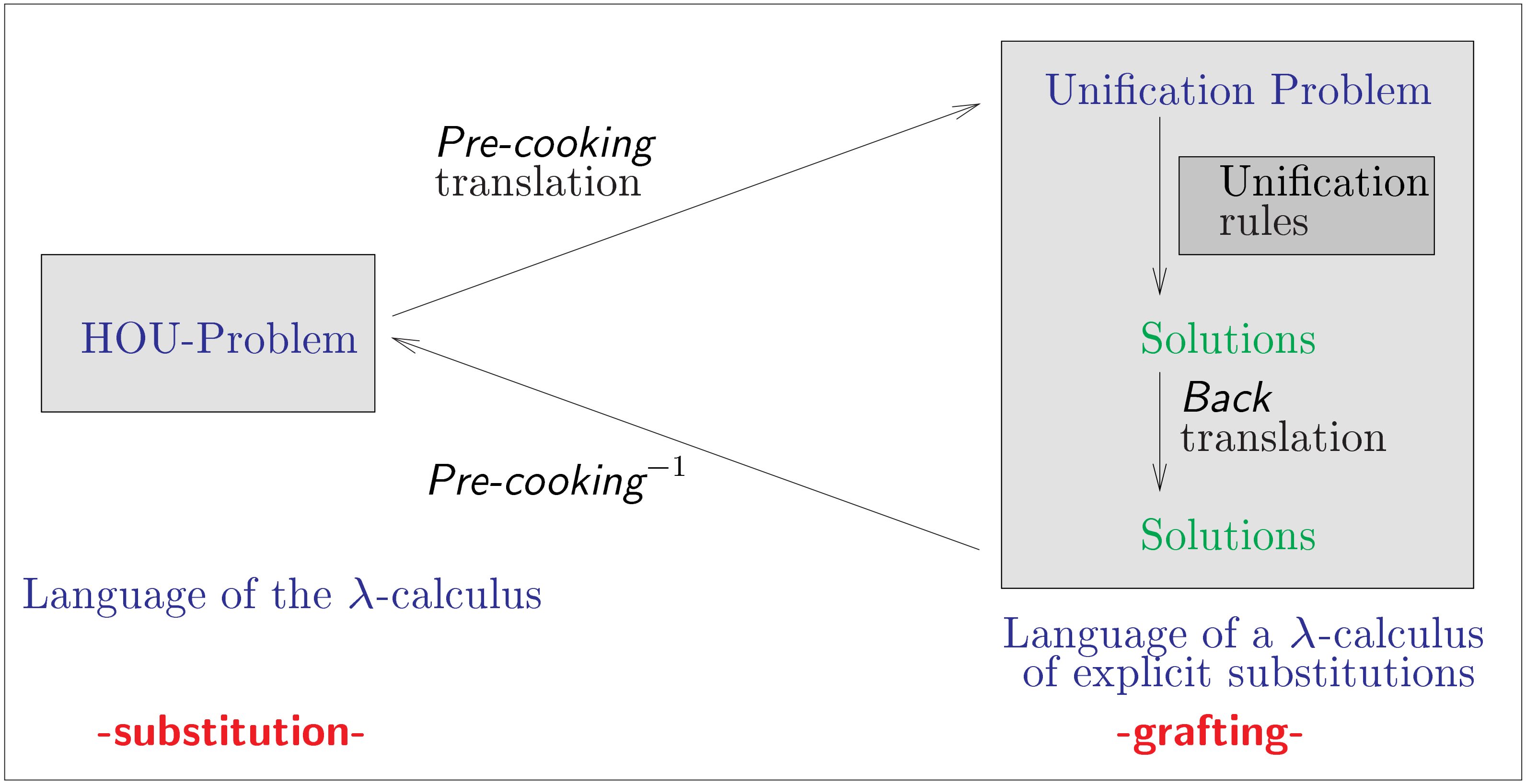
- Higher order *substitution*: $\{X/1\}(\lambda.(1\ \ X)\ \ X) = (\lambda.(1\ \ 2)\ \ 1)$

substitution	$\neq$	*grafting*
$\{X/a\}(\lambda.X)$		$(\lambda.X)\{X/a\}$
$\Vert$		$\Vert$
$\lambda.\{X/a^+\}X$		$\lambda.X\{X/a\}$
$\Vert$		$\Vert$
$\lambda.\underbrace{a^+}_{\text{lift}}$	$\neq$	$\lambda.a$

β-reduction

$$(\lambda.a\ \ b) \to \{1/b\}a$$

General HOU method via explicit substitutions



- Introduced by G. Dowek, T. Hardin and C. Kirchner using the $\lambda\sigma$-calculus.
- Subsumes Huet's HOU method.

2. HOU in the λs_e-calculus

2.1. Unification in the λs_e-style of explicit substitutions

2.2. Checking arithmetic constraints (versus shifts and composition in $\lambda\sigma$)

2.3. Translations between the pure λ-calculus and the λs_e-calculus

2.1. Unification in the λs_e-style of explicit substitution

- Terms in λs_e: $a ::= \mathcal{X} \mid \mathbb{N} \mid (a\ \ a) \mid \lambda.a \mid a\sigma^j a \mid \varphi_k^i a$, for $j,\ i \geq 1\,,\ \ k \geq 0$ where $\mathcal{X}$ *meta-variables* and $\mathbb{N}$ set of de Bruijn indices.

- A λs_e-**unification problem** P is: $\left\{ \bigvee_{j \in J} \underbrace{\exists \vec{w}_j \bigwedge_{i \in I_j} s_i =^?_{\lambda s_e} t_i}_{\text{unification system}} \right.$

- A **unifier** of $\underbrace{\exists \vec{w} \bigwedge_{i \in I} s_i =^?_{\lambda s_e} t_i}_{\text{unification system}}$ is a **grafting** σ such that $\boxed{\exists \vec{w} \bigwedge_{i \in I} s_i\sigma = t_i\sigma}$

Example : $(\lambda.(\lambda.(X\ \ 2)\ \ 1)\ \ Y) =^?_{\lambda s_e} (\lambda.(Z\ \ 1)\ \ U)$

$\downarrow$ $X, Z : A \rightarrow A;\ Y, U : A$

Normalize $((X\sigma^2 Y)\sigma^1(\varphi^1_0 Y)\ \ \varphi^1_0 Y) =^?_{\lambda s_e} (Z\sigma^1 U\ \ \varphi^1_0 U)$

$\downarrow$

Dec-App $(X\sigma^2 Y)\sigma^1(\varphi^1_0 Y) =^?_{\lambda s_e} Z\sigma^1 U \ \wedge\ \varphi^1_0 Y =^?_{\lambda s_e} \varphi^1_0 U$

$\downarrow$

Dec-φ $(X\sigma^2 Y)\sigma^1(\varphi^1_0 Y) =^?_{\lambda s_e} Z\sigma^1 U \ \wedge\ Y =^?_{\lambda s_e} U$

$\downarrow$

Replace $(X\sigma^2 Y)\sigma^1(\varphi^1_0 Y) =^?_{\lambda s_e} Z\sigma^1 Y \ \wedge\ Y =^?_{\lambda s_e} U$

$\downarrow_*$

Exp-λ + Replace

$$((\lambda.X')\sigma^2 Y)\sigma^1(\varphi^1_0 Y) =^?_{\lambda s_e} (\lambda.Z')\sigma^1 Y \ \wedge\ \begin{cases} Y =^?_{\lambda s_e} U \\ X =^?_{\lambda s_e} \lambda.X' \\ Z =^?_{\lambda s_e} \lambda.Z' \end{cases}$$

$\downarrow_*$

Normalize + Dec-λ

$$(X'\sigma^3 Y)\sigma^2(\varphi^1_0 Y) =^?_{\lambda s_e} Z'\sigma^2 Y \ \wedge\ \begin{cases} Y =^?_{\lambda s_e} U \\ X =^?_{\lambda s_e} \lambda.X' \\ Z =^?_{\lambda s_e} \lambda.Z' \end{cases}$$

- *Solved* equations: $\left\{\begin{array}{l} Y =^?_{\lambda s_e} U \\ X =^?_{\lambda s_e} \lambda.X' \\ Z =^?_{\lambda s_e} \lambda.Z' \end{array}\right.$

- *Flex-Flex* equations: $(X'\sigma^3 Y)\sigma^2(\varphi^1_0 Y) =^?_{\lambda s_e} Z'\sigma^2 Y$

(both groups together:) *Solved Forms*

• Solutions: $\{Y/X_1, U/X_1\} \bigcup$ solutions for X and Z given by the *Flex-Flex* equation.

Take, for instance, $\{Y/X_1, U/X_1\} \bigcup \{X/\lambda.\mathtt{n}+1, Z/\lambda.\mathtt{n}\}$ with $n > 2$:
$\underline{(\lambda.(\lambda.(\lambda.\mathtt{n}+1\ \ 2)\ \ 1)\ \ X_1)} \rightarrow_\beta (\lambda.(\lambda.\mathtt{n}\ \ 2)\ \ X_1) \rightarrow_\beta (\lambda.\mathtt{n}-1\ \ X_1) \rightarrow_\beta \underline{\mathtt{n}-2}$
and
$\underline{(\lambda.(\lambda.\mathtt{n}\ \ 1)\ \ X_1)} \rightarrow_\beta (\lambda.\mathtt{n}-1\ \ X_1) \rightarrow_\beta \underline{\mathtt{n}-2}$

- Correctness: If P reduces to P' then every unifier of P' is a unifier of P.
- Completeness: If P reduces to P' then every unifier of P is a unifier of P'.

Theorem [Correctness and Completeness]

The λs_e-unification rules are correct and complete.

2.2. Checking arithmetic constraints (versus shifts and composition in $\lambda\sigma$)

λs_e-calculus and λ-calculus $\rightarrow$ *Term*, *Substitution* } objects $\lambda\sigma$-calculus

λs_e uses all de Bruijn indices: $\mathbb{N}$

$\lambda\sigma$ uses only $\mathtt{1}$, "shift" and "composition": $\mathtt{n} \equiv \mathtt{1}[\underbrace{\uparrow \circ \cdots \circ \uparrow}_{n-1}]$

2.3. Translations between the pure λ-calculus and the λs_e-calculus

- A unifier of $\lambda.X =_{\beta\eta} \lambda.a$ is not a $\{X/b\}$ such that $b =_{\beta\eta} a$:

$$\{X/b\}(\lambda.X) = \lambda.(\{X/b^+\}X) = \lambda.(X\{X/b^+\}) = \lambda.b^+$$

- The **pre-cooking** of a λ-term in de Bruijn notation into the λs_e-calculus is defined by $\boldsymbol{a_{pc}} = PC(a, 0)$ where $PC(a, n)$ is defined by:

1. $PC(\lambda_B.a, n) = \lambda_B.PC(a, n+1)$
2. $PC((a\ \ b), n) = (PC(a, n)\ \ PC(b, n))$
3. $PC(\mathtt{k}, n) = \mathtt{k}$
4. $PC(X, n) = \begin{cases} \text{if } n = 0 \text{ then } X \\ \text{else } \varphi_0^{n+1} X \end{cases}$

Proposition[Semantics of pre-cooking]

$$\underbrace{(\{X_1/b_1, \ldots, X_p/b_p\}(a))_{pc}}_{\text{Substitution}} = \underbrace{a_{pc}\{X_1/b_{1_{pc}}, \ldots, X_p/b_{p_{pc}}\}}_{\text{Grafting}}$$

Proposition[Correspondence between solutions]

$$\exists N_1,\ldots,N_p \quad \underbrace{\{X_1/N_1,\ldots,X_p/N_p\}}_{\text{substitution}}(a) =_{\beta\eta} \underbrace{\{X_1/N_1,\ldots,X_p/N_p\}}_{\text{substitution}}(b)$$

$$\Longleftrightarrow$$

$$\exists M_1,\ldots,M_p \quad a_{pc}\underbrace{\{X_1/M_1,\ldots,X_p/M_p\}}_{\text{grafting}} =_{\lambda s_e} b_{pc}\underbrace{\{X_1/M_1,\ldots,X_p/M_p\}}_{\text{grafting}}$$

3. A simple example

Problem: $\lambda.(X\ \ 2) =^?_{\beta\eta} \lambda.2, \quad 2 : A, \quad X : A \rightarrow A$

Solution: Pre-cooking and then as before:

$$\lambda.(\varphi^2_0(X)\ \ 2) =^?_{\lambda s_e} \lambda.2 \qquad \rightarrow \textit{Dec-}\lambda$$

$$(\varphi^2_0(X)\ \ 2) =^?_{\lambda s_e} 2 \qquad \rightarrow \textit{Exp-}\lambda$$

$$\exists Y\, (\varphi^2_0(X)\ \ 2) =^?_{\lambda s_e} 2 \wedge X =^?_{\lambda s_e} \lambda.Y \qquad \rightarrow \textit{Replace}$$

$$\exists Y\, (\varphi^2_0(\lambda.Y)\ \ 2) =^?_{\lambda s_e} 2 \wedge X =^?_{\lambda s_e} \lambda.Y \qquad \rightarrow \textit{Normalize}$$

$$\exists Y\, (\varphi^2_1 Y)\sigma^1 2 =^?_{\lambda s_e} 2 \wedge X =^?_{\lambda s_e} \lambda.Y \qquad \rightarrow \textit{Exp-app}$$

$$(\exists Y\, (\varphi^2_1 Y)\sigma^1 2 =^?_{\lambda s_e} 2 \wedge X =^?_{\lambda s_e} \lambda.Y) \wedge (Y =^?_{\lambda s_e} 1 \vee Y =^?_{\lambda s_e} 2) \qquad \rightarrow \textit{Replace}$$

$$((\varphi^2_1 1)\sigma^1 2 =^?_{\lambda s_e} 2 \wedge X =^?_{\lambda s_e} \lambda.1) \vee ((\varphi^2_1 2)\sigma^1 2 =^?_{\lambda s_e} 2 \wedge X =^?_{\lambda s_e} \lambda.2) \qquad \rightarrow \textit{Normalize}$$

$$(2 =^?_{\lambda s_e} 2 \wedge X =^?_{\lambda s_e} \lambda.1) \vee (2 =^?_{\lambda s_e} 2 \wedge X =^?_{\lambda s_e} \lambda.2) \qquad \equiv$$

$$(X =^?_{\lambda s_e} \lambda.1) \vee (X =^?_{\lambda s_e} \lambda.2)$$

Problem: $\boxed{\lambda.(X\ \ 2) =^?_{\beta\eta} \lambda.2, \quad 2 : A, \quad X : A \rightarrow A}$

Solutions: $\begin{cases} \{X/\lambda.1\} \\ \{X/\lambda.2\} \end{cases}$

Note that we have:

$$\{X/\lambda.1\}(\lambda.(X\ \ 2)) = \lambda.(\{X/(\lambda.1)^+\}(X)\ \ 2) = \\ \lambda.(\lambda.1^{+1}\ \ 2) = \lambda.(\lambda.1\ \ 2) =_\beta \lambda.2$$

and

$$\{X/\lambda.2\}(\lambda.(X\ \ 2)) = \lambda.(\{X/(\lambda.2)^+\}(X)\ \ 2) = \\ \lambda.(\lambda.2^{+1}\ \ 2) = \lambda.(\lambda.3\ \ 2) =_\beta \lambda.2$$

4. HOU via explicit substituions in the praxis

Types were omitted!

Typing rules for the lambda calculus in *de Bruijn* notation:

$$\frac{1 \leq i \leq n}{A_1.A_2.\ldots.A_n \vdash \mathtt{i} : A_i} \text{ (Var)}$$

$$\frac{A.\Gamma \vdash M : B}{\Gamma \vdash \lambda_A.M : A \to B} \text{ (Abs)}$$

$$\frac{\Gamma \vdash M : A \to B \quad \Gamma \vdash N : A}{\Gamma \vdash (M \;\; N) : B} \text{ (Appl)}$$

Table 1: Undecorated and decorated typing rules for the λs_e-calculus

(Var)	$A.\Gamma \vdash \mathtt{1} : A$	$\mathtt{1}_A^{A.\Gamma}$
(Varn)	$\dfrac{\Gamma \vdash \mathtt{n} : B}{A.\Gamma \vdash \mathtt{n}+\mathtt{1} : B}$	$\dfrac{\mathtt{n}_B^{\Gamma}}{(\mathtt{n}+\mathtt{1})_B^{A.\Gamma}}$
(Lambda)	$\dfrac{A.\Gamma \vdash b : B}{\Gamma \vdash \lambda_A.b : A \to B}$	$\dfrac{b_B^{A.\Gamma}}{(\lambda_A.b_B^{A.\Gamma})_{A\to B}^{\Gamma}}$
(App)	$\dfrac{\Gamma \vdash b : A \to B \qquad \Gamma \vdash a : A}{\Gamma \vdash (b\ \ a) : B}$	$\dfrac{b_{A\to B}^{\Gamma}\ ,\ a_A^{\Gamma}}{(b_{A\to B}^{\Gamma}\ \ a_A^{\Gamma})_B^{\Gamma}}$
(Sigma)	$\dfrac{\Gamma_{\geq i} \vdash b : B \qquad \Gamma_{<i}.B.\Gamma_{\geq i} \vdash a : A}{\Gamma \vdash a\,\sigma^i b : A}$	$\dfrac{b_B^{\Gamma_{\geq i}}\ ,\ a_A^{\Gamma_{<i}.B.\Gamma_{\geq i}}}{(a_A^{\Gamma_{<i}.B.\Gamma_{\geq i}}\ \sigma^i b_B^{\Gamma_{\geq i}})_A^{\Gamma}}$
(Phi)	$\dfrac{\Gamma_{\leq k}.\Gamma_{\geq k+i} \vdash a : A}{\Gamma \vdash \varphi_k^i a : A}$	$\dfrac{a_A^{\Gamma_{\leq k}.\Gamma_{\geq k+i}}}{(\varphi_k^i a_A^{\Gamma_{\leq k}.\Gamma_{\geq k+i}})_A^{\Gamma}}$
(Meta)	$\Gamma_X \vdash X : A_X$	$X_{A_X}^{\Gamma_X}$

Decoration of *substitution objects* in the $\lambda\sigma$-calculus: $s \triangleright \Gamma$.
Types are environments!

Decorated typing rules in the $\lambda\sigma$-calculus:

(Shift)	$\dfrac{}{\uparrow^{A.\Gamma}_{\Gamma}}$
(Comp)	$\dfrac{s^{\Theta}_{\Gamma} \ , \ t^{\Delta}_{\Theta}}{(s^{\Theta}_{\Gamma} \circ t^{\Delta}_{\Theta})^{\Delta}_{\Gamma}}$
(Clos)	$\dfrac{a^{\Delta}_{A} \ , \ s^{\Gamma}_{\Delta}}{(a^{\Delta}_{A}[s^{\Gamma}_{\Delta}])^{\Gamma}_{A}}$

The de Bruijn index n is decorated in linear time and space in λs_e while its corresponding $\lambda\sigma$-term $1[\uparrow^{n-1}]$ in quadratic space and time:

$$
(comp)\dfrac{(shift)\ \uparrow^{A_{n-1}.A_n.\Gamma}_{A_n.\Gamma} \quad , \quad (shift)\ \uparrow^{A_{n-2}\ldots\Gamma}_{A_{n-1}.A_n.\Gamma}}{(\uparrow^{A_{n-1}.A_n.\Gamma}_{A_n.\Gamma} \circ \uparrow^{A_{n-2}\ldots\Gamma}_{A_{n-1}.A_n.\Gamma})^{A_{n-2}\ldots\Gamma}_{A_n.\Gamma} \quad , \quad (shift)\ \uparrow^{A_{n-3}\ldots A_n.\Gamma}_{A_{n-2}\ldots\Gamma}}
$$

$$
(comp)\dfrac{(\uparrow^{A_{n-1}.A_n.\Gamma}_{A_n.\Gamma} \circ \uparrow^{A_{n-2}\ldots\Gamma}_{A_{n-1}.A_n.\Gamma})^{A_{n-2}\ldots\Gamma}_{A_n.\Gamma} \quad , \quad (shift)\ \uparrow^{A_{n-3}\ldots A_n.\Gamma}_{A_{n-2}\ldots\Gamma}}{\vdots}
$$

$$
(comp)\dfrac{\vdots}{(\cdots(\uparrow^{A_{n-1}.A_n.\Gamma}_{A_n.\Gamma} \circ \uparrow^{A_{n-2}\ldots\Gamma}_{A_{n-1}.A_n.\Gamma})^{A_{n-2}\ldots\Gamma}_{A_n.\Gamma} \circ \cdots)^{A_1\ldots A_n.\Gamma}_{A_n.\Gamma} \quad , \quad (var) 1^{A_n.\Gamma}_{A_n}}
$$

$$
(clos)\dfrac{(\cdots(\uparrow^{A_{n-1}.A_n.\Gamma}_{A_n.\Gamma} \circ \uparrow^{A_{n-2}\ldots\Gamma}_{A_{n-1}.A_n.\Gamma})^{A_{n-2}\ldots\Gamma}_{A_n.\Gamma} \circ \cdots)^{A_1\ldots A_n.\Gamma}_{A_n.\Gamma} \quad , \quad (var) 1^{A_n.\Gamma}_{A_n}}{(1^{A_n.\Gamma}_{A_n}[(\cdots(\uparrow^{A_{n-1}.A_n.\Gamma}_{A_n.\Gamma} \circ \uparrow^{A_{n-2}\ldots\Gamma}_{A_{n-1}.A_n.\Gamma})^{A_{n-2}\ldots\Gamma}_{A_n.\Gamma} \circ \cdots)^{A_1\ldots A_n.\Gamma}_{A_n.\Gamma}])^{A_1\ldots A_n.\Gamma}_{A_n}}
$$

Lemma [Linear versus quadratic decorations] Pre-cooked λ-terms in the λs_e-calculus have linear decorations on the size of the λ-terms and the magnitude of their de Bruijn indices, while in $\lambda\sigma$ these decorations are quadratic.

- Additionally, some rules are *expansive* in $\lambda\sigma$:

Consider the decorated λ-term: $((\lambda_A.((\lambda_A.X_A^{A.A.A.\Gamma})_{A\to A}^{A.A.\Gamma}\ \ 1_A^{A.A.\Gamma})_A^{A.A.\Gamma})_{A\to A}^{A.\Gamma}\ \ 1_A^{A.\Gamma})_A^{A.\Gamma}$
(i.e.,$(\lambda_A.(\lambda_A.X\ \ 1)\ \ 1))$

- Applying the *Beta* rule in λs_e: $\begin{cases} \to_{Beta} & ((\lambda_A.(X_A^{A.A.A.\Gamma}\sigma^1 1_A^{A.A.\Gamma})_A^{A.A.\Gamma})_{A\to A}^{A.\Gamma}\ \ 1_A^{A.\Gamma})_A^{A.\Gamma} \\ \to_{Beta} & ((X_A^{A.A.A.\Gamma}\sigma^1 1_A^{A.A.\Gamma})_A^{A.A.\Gamma}\sigma^1 1_A^{A.\Gamma})_A^{A.\Gamma} \end{cases}$

- In $\lambda\sigma$: $\begin{cases} \to_{Beta} & ((\lambda_A.(X_A^{A.A.A.\Gamma}[(1_A^{A.A.\Gamma}.id_{A.A.\Gamma}^{A.A.\Gamma})_{A.A.A.\Gamma}^{A.A.\Gamma}])_A^{A.A.\Gamma})_{A\to A}^{A.\Gamma}\ \ 1_A^{A.\Gamma})_A^{A.\Gamma} \\ \to_{Beta} & ((X_A^{A.A.A.\Gamma}[(1_A^{A.A.\Gamma}.id_{A.A.\Gamma}^{A.A.\Gamma})_{A.A.A.\Gamma}^{A.A.\Gamma}])_A^{A.A.\Gamma}[(1_A^{A.\Gamma}.id_{A.\Gamma}^{A.\Gamma})_{A.A.\Gamma}^{A.\Gamma}])_A^{A.\Gamma} \end{cases}$

5. Related work

Our development of the λs_e-HOU was based on the original ones of Dowek, Hardin and Kirchner for the $\lambda\sigma$-calculus of explicit substitutions.

One of our motivations was, in the practical setting of HOU, to compare the advantages and disadvantages of the two styles of explicit substitutions. This provides objective facts about that interesting theoretical question.

We think that our method can be adapted for applications in/for systems as the λProlog, Maude and ELAN.

6. Future work and Conclusions

To be done:
- Prototype implementation.
- Comparison with the *suspension* calculus.

- $\lambda\sigma$-(HO)Unification and λs_e-(HO)Unification strategies don't differ.
- Pre-cooking (and back) translations in $\lambda\sigma$ and λs_e differ:
  - A simple selection of the scripts for the operators φ and σ in λs_e corresponds to the manipulation of substitution objects in the $\lambda\sigma$-HOU approach.
  - Use of all de Bruijn indices makes our approach simpler.

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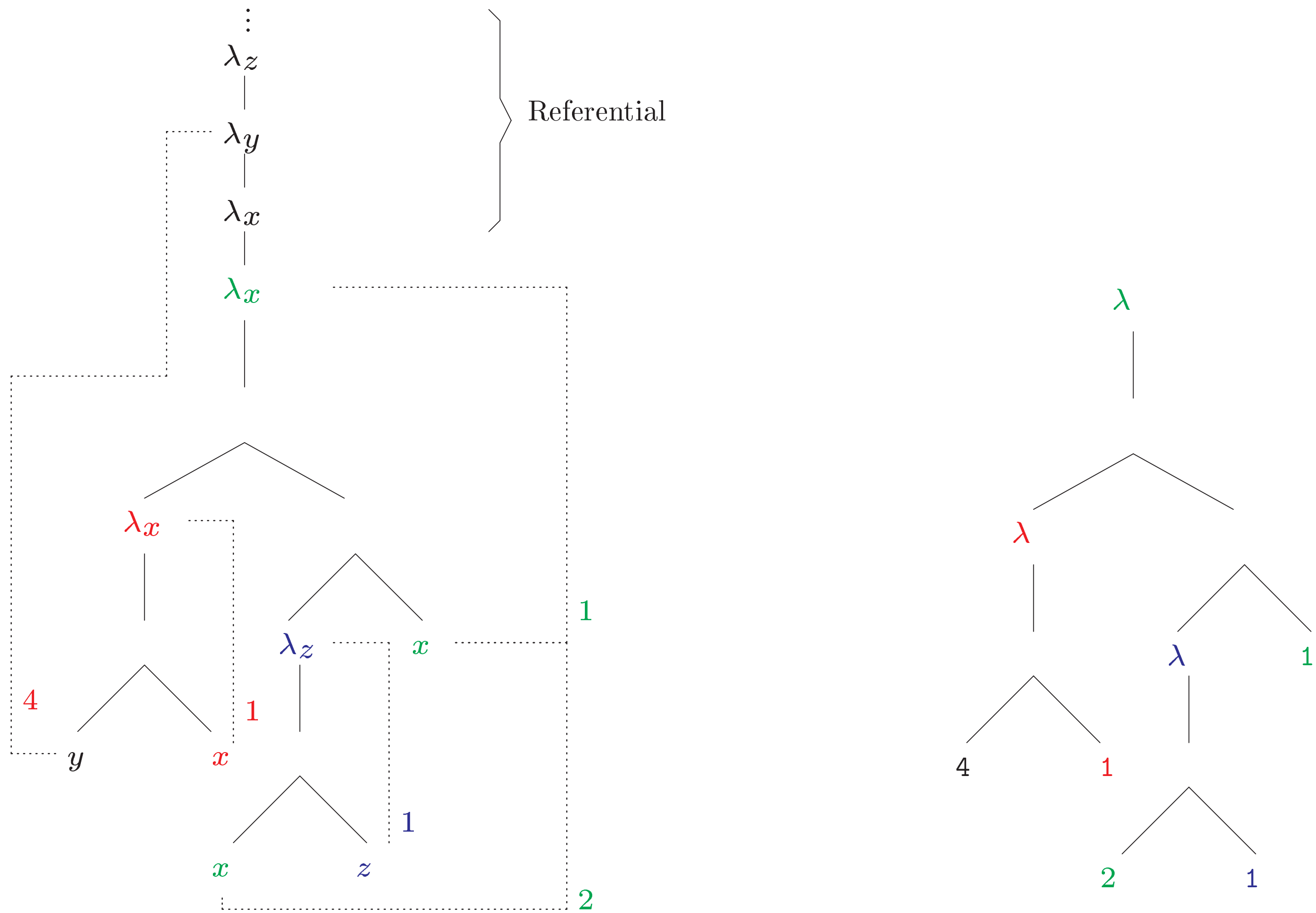
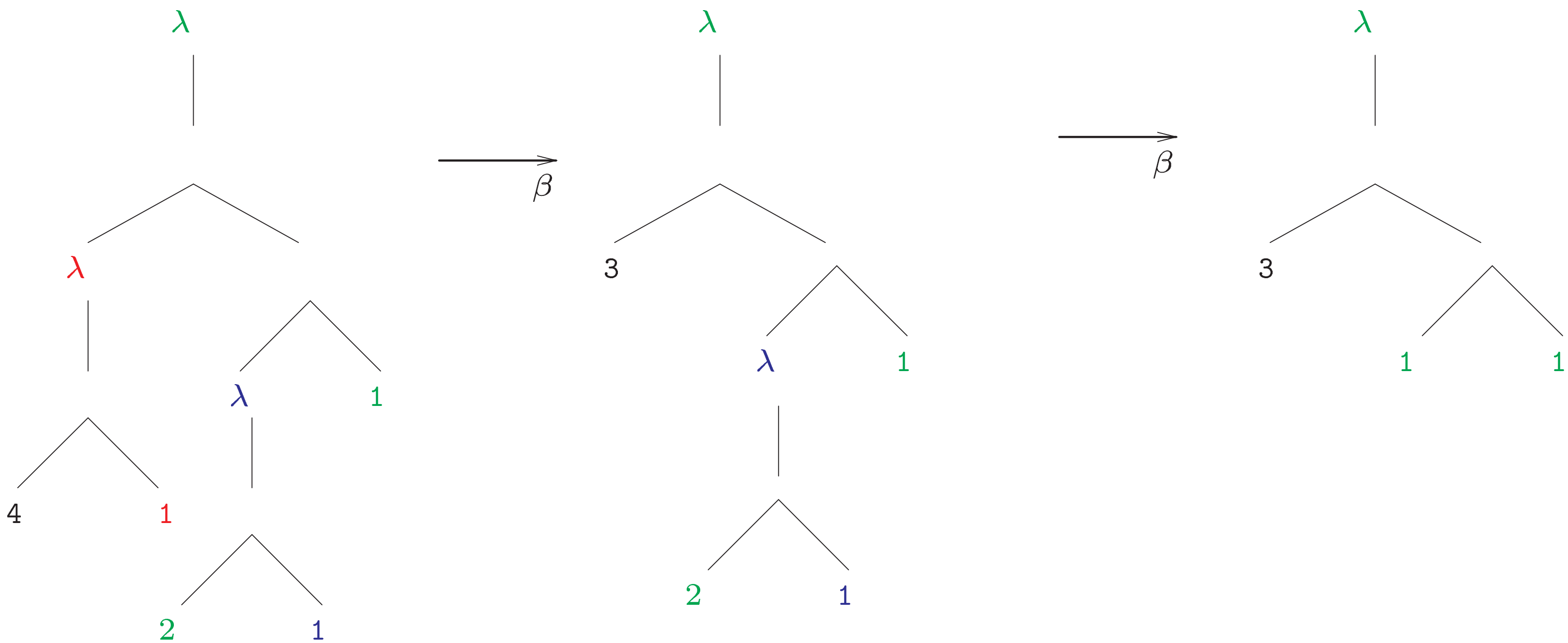


Figure 1: $\lambda_x.(\lambda_x.(y\ \ x)\ \ (\lambda_z.(x\ \ z)\ \ x))$ and its de Bruijn version: $\lambda.(\lambda.(4\ \ 1)\ \ (\lambda.(2\ \ 1)\ \ 1))$


λ
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1

Exp-App $\lambda\sigma$-unification rule

$$P \wedge X[a_1 \ldots a_p . \uparrow^n] =^?_{\lambda\sigma} (\mathtt{m}\, b_1 \ldots b_q) \quad \rightarrow$$

$$\bigwedge \left\{ \begin{array}{l} P \\ X[a_1 \ldots a_p . \uparrow^n] =^?_{\lambda\sigma} (\mathtt{m}\, b_1 \ldots b_q) \\ \bigvee_{r \in R_p \cup R_i} \exists H_1 \ldots H_k, X =^?_{\lambda\sigma} (\mathtt{r}\ H_1 \ldots H_k) \end{array} \right.$$

X not solved and atomic; $H_1, \ldots, H_k$ variables of appropriate types; $\Gamma_{H_i} = \Gamma_X$, $R_p \subseteq \{1, \ldots, p\}$ such that $(\mathtt{r}\ H_1 \ldots H_k)$ has the right type, $R_i =$ if $m \geq n+1$ then $\{m - n + p\}$ else $\emptyset$

Exp-App λs_e-unification rule

$$P \wedge \psi_{i_p}^{j_p} \ldots \psi_{i_1}^{j_1}(X, a_1, \ldots, a_p) =^?_{\lambda s_e} (\mathtt{m}\; b_1 \ldots b_q) \quad \rightarrow$$

$$\bigwedge \begin{cases} P \\ \psi_{i_p}^{j_p} \ldots \psi_{i_1}^{j_1}(X, a_1, \ldots, a_p) =^?_{\lambda s_e} (\mathtt{m}\; b_1 \ldots b_q) \\ \bigvee_{r \in R_p \cup R_i} \exists H_1, \ldots, H_k, X =^?_{\lambda s_e} (\mathtt{r}\; H_1 \ldots H_k) \end{cases}$$

$\psi_{i_p}^{j_p} \ldots \psi_{i_1}^{j_1}(X, a_1, \ldots, a_p)$ skeleton of a λs_e-normal term; X atomic and not solved; $\Gamma_{H_i} = \Gamma_X$, $R_p \subseteq \{i_1, \ldots, i_p\}$ of superscripts of the σ operator such that $(\mathtt{r}\; H_1 \ldots H_k)$ has the right type, $R_i = \bigcup_{k=0}^{p}$ if $i_k \geq m + p - k - \sum_{l=k+1}^{p} j_l > i_{k+1}$ then $\{m + p - k - \sum_{l=k+1}^{p} j_l\}$ else $\emptyset$, where $i_0 = \infty, i_{p+1} = 0$

In the $\lambda\sigma$-calculus

$X[a_1 \ldots a_p. \uparrow^n] =^?_{\lambda\sigma} (\mathtt{m}\, b_1 \ldots b_q)$ has solutions of the form:

$$\left(\mathtt{1}[\underbrace{\uparrow \circ \cdots \circ \uparrow}_{r-1}] \quad \underbrace{H_1 \quad \ldots \quad H_k}_{\text{of appropriate type}} \right)$$

$$\mathtt{1}[\underbrace{\uparrow \circ \cdots \circ \uparrow}_{r-1}] \quad [a_1 \ldots a_p. \uparrow^n] \quad = \begin{cases} a_i, \text{ if } 1 \leq r = i \leq p \\ \mathtt{1}[\underbrace{\uparrow \circ \cdots \circ \uparrow}_{r-1-p}] \quad [\underbrace{\uparrow \circ \cdots \circ \uparrow}_{n}] \quad otherwise. \end{cases}$$

In the λs_e-calculus

$\psi_{k_p}^{j_p} \ldots \psi_{k_1}^{j_1}(X, a_1, \ldots, a_p) =^?_{\lambda s_e} (\mathtt{m}\ b_1 \ldots b_q)$ solutions of the form:

$$\left(\mathtt{n}\ \underbrace{H_1 \quad \ldots \quad H_k}_{\text{of appropriate type}}\right)$$

such that for some i, $\left[\begin{array}{c} k_{i+1} < n \leq k_i \\ \text{and} \\ n - (p - i) + \sum_{r=i+1}^{p} j_r = m \end{array}\right]$

$(\lambda M \ \ N)$	$\xrightarrow{(Beta)}$	$M[N \cdot id]$	$4(0) \to 5(2)$
$(M \ \ N)[S]$	$\xrightarrow{(App)}$	$(M[S] \ \ N[S])$	$5(1) \to 7(2)$
$(\lambda M)[S]$	$\xrightarrow{(Abs)}$	$\lambda M[1 \cdot (S \circ \uparrow)]$	$4(1) \to 8(4)$
$M[S][T]$	$\xrightarrow{(Clos)}$	$M[S \circ T]$	$5(2) \to 5(3)$
$1[M \cdot S]$	$\xrightarrow{(VarCons)}$	M	$5(2) \to 1(0)$
$M[id]$	$\xrightarrow{(Id)}$	M	$3(1) \to 1(0)$
$(S_1 \circ S_2) \circ T$	$\xrightarrow{(Assoc)}$	$S_1 \circ (S_2 \circ T)$	$5(5) \to 5(5)$
$(M \cdot S) \circ T$	$\xrightarrow{(Map)}$	$M[T] \cdot (S \circ T)$	$5(4) \to 7(5)$
$id \circ S$	$\xrightarrow{(IdL)}$	S	$3(3) \to 1(1)$
$S \circ id$	$\xrightarrow{(IdR)}$	S	$3(3) \to 1(1)$
$\uparrow \circ (M \cdot S)$	$\xrightarrow{(ShiftCons)}$	S	$5(4) \to 1(1)$
$1 \cdot \uparrow$	$\xrightarrow{(VarShift)}$	id	$3(2) \to 1(1)$
$1[S] \cdot (\uparrow \circ S)$	$\xrightarrow{(SCons)}$	S	$7(5) \to 1(1)$
$\lambda(M \ \ 1)$	$\xrightarrow{(Eta)}$	N if $M =_\sigma N[\uparrow]$	$4(0) \to 4(1)$

Table 2: Rewriting rules of the $\lambda\sigma$-calculus

$(\lambda M \ \ N)$	$\xrightarrow{(\sigma\text{-generation})}$	$M\,\sigma^1\,N$	$4 \to 3$
$(\lambda M)\,\sigma^i N$	$\xrightarrow{(\sigma\text{-}\lambda\text{-transition})}$	$\lambda(M\,\sigma^{i+1}\,N)$	$4 \to 4$
$(M_1 \ \ M_2)\,\sigma^i N$	$\xrightarrow{(\sigma\text{-app-transition})}$	$((M_1\,\sigma^i N) \ \ (M_2\,\sigma^i N))$	$5 \to 7$
$\mathtt{n}\,\sigma^i N$	$\xrightarrow{(\sigma\text{-destruction})}$	$\begin{cases} \mathtt{n-1} & \text{if } n > i \\ \varphi^i_0\,N & \text{if } n = i \\ \mathtt{n} & \text{if } n < i \end{cases}$	$3 \to 1, 2, 1$
$\varphi^i_k(\lambda M)$	$\xrightarrow{(\varphi\text{-}\lambda\text{-transition})}$	$\lambda(\varphi^i_{k+1}\,M)$	$3 \to 3$
$\varphi^i_k(M_1 \ \ M_2)$	$\xrightarrow{(\varphi\text{-app-transition})}$	$((\varphi^i_k\,M_1) \ \ (\varphi^i_k\,M_2))$	$4 \to 5$
$\varphi^i_k\,\mathtt{n}$	$\xrightarrow{(\varphi\text{-destruction})}$	$\begin{cases} \mathtt{n+i-1} & \text{if } n > k \\ \mathtt{n} & \text{if } n \leq k \end{cases}$	$2 \to 1, 1$
$(M_1\,\sigma^i M_2)\,\sigma^j\,N$	$\xrightarrow{(\sigma\text{-}\sigma\text{-transition})}$	$(M_1\,\sigma^{j+1}\,N)\,\sigma^i\,(M_2\,\sigma^{j-i+1}\,N)$ if $i \leq j$	$5 \to 7$
$(\varphi^i_k\,M)\,\sigma^j\,N$	$\xrightarrow{(\sigma\text{-}\varphi\text{-transition 1})}$	$\varphi^{i-1}_k\,M$ if $k < j < k+i$	$4 \to 2$
$(\varphi^i_k\,M)\,\sigma^j\,N$	$\xrightarrow{(\sigma\text{-}\varphi\text{-transition 2})}$	$\varphi^i_k(M\,\sigma^{j-i+1}\,N)$ if $k+i \leq j$	$4 \to 4$
$\varphi^i_k(M\,\sigma^j\,N)$	$\xrightarrow{(\varphi\text{-}\sigma\text{-transition})}$	$(\varphi^i_{k+1}\,M)\,\sigma^j\,(\varphi^i_{k+1-j}\,N)$ if $j \leq k+1$	$4 \to 5$
$\varphi^i_k\,(\varphi^j_l\,M)$	$\xrightarrow{(\varphi\text{-}\varphi\text{-transition 1})}$	$\varphi^j_l\,(\varphi^i_{k+1-j}\,M)$ if $l+j \leq k$	$3 \to 3$
$\varphi^i_k\,(\varphi^j_l\,M)$	$\xrightarrow{(\varphi\text{-}\varphi\text{-transition 2})}$	$\varphi^{j+i-1}_l\,M$ if $l \leq k < l+j$	$3 \to 2$
$\lambda(M \ \ \mathtt{1})$	$\xrightarrow{(\text{Eta})}$	N if $M =_{s_e} \varphi^2_0 N$	$4 \to 3$

Table 3: Rewriting rules of the λs_e-calculus