

Algebraic Properties of Word Equations

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- Historical excursus
 - Fine and Wilf theorem (two original but less known proofs)
 - Compactness property of systems of word equations
- Encoding of word equations into polynomials
[including a commercial break]

Fine and Wilf: The Periodicity Lemma, Proof 1

Theorem

Let $\{f_n\}_0^\infty$ and $\{g_n\}_0^\infty$ be two periodic sequences of periods h and k , respectively. If $f_n = g_n$ for $h + k - (h, k)$ consecutive integers n , then $f_n = g_n$ for all n . The result would be false if $h + k - (h, k)$ were replaced by anything smaller.

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Proof.

Use discrete Fourier bases of functions with the period k and h :

$$\Psi : \quad \psi_j(n) = e^{2\pi i \frac{j}{k} n}, \quad j = 0, 1, \dots, k-1,$$

$$\Phi : \quad \phi_j(n) = e^{2\pi i \frac{j}{h} n}, \quad j = 0, 1, \dots, h-1.$$

There are altogether $h + k - (h, k)$ linearly independent elements in $\Psi \cup \Phi$. Therefore the interpolation of the shared interval is unique iff it is of at least that length (with the support in $\Psi \cap \Phi$). \square

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Proof.

$$F(x) = \sum_0^\infty f_n x^n = \frac{P(x)}{1 - x^h} \qquad G(x) = \sum_0^\infty g_n x^n = \frac{Q(x)}{1 - x^k}$$

$$F(x) - G(x) = \frac{P(x)(1 - x^k)/(1 - x^{(h,k)}) - Q(x)(1 - x^h)/(1 - x^{(h,k)})}{(1 - x^h)(1 - x^k)/(1 - x^{(h,k)})}$$



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Easily equivalent (1980) to “Ehrenfeucht’s conjecture”
(beginning of 1970s - Nowa Księga Szkocka, problem 105)

- **Theorem**

Every language over a finite alphabet has a finite test set (testing equality of morphisms on the language).

- Proved independently by Albert & Lawrence (1985); and Guba (1986).
- Core of both proofs: Hilbert’s basis theorem.

$$\mathbf{a} = \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}, \quad \mathbf{b} = \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix} \quad \mathrm{SL}_2(\mathbb{N}_0) = \langle \mathbf{a}, \mathbf{b} \rangle \cong \{a, b\}^*$$

$$\mathbf{c} = \begin{pmatrix} 1 & 2 \\ 0 & 1 \end{pmatrix}, \quad \mathbf{d} = \begin{pmatrix} 1 & 0 \\ 2 & 1 \end{pmatrix} \quad \langle \mathbf{c}, \mathbf{d} \rangle \cong F_2$$

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$$x_j \mapsto m_j = \begin{pmatrix} a_j & b_j \\ c_j & d_j \end{pmatrix}$$

$$\begin{array}{ccc} \Theta^* & \xrightarrow{\psi} & M \\ & \searrow \varphi & \swarrow \tilde{\varphi} \\ & \mathrm{SL}_2(\mathbb{N}_0) & \end{array}$$

$$xyz = zyx$$

$$x \mapsto \begin{pmatrix} a_x & b_x \\ c_x & d_x \end{pmatrix} \quad y \mapsto \begin{pmatrix} a_y & b_y \\ c_y & d_y \end{pmatrix} \quad z \mapsto \begin{pmatrix} a_z & b_z \\ c_z & d_z \end{pmatrix}$$

$$a_z b_x c_y + a_x b_y c_z + b_x c_z d_y = a_z b_y c_x + a_x b_z c_y + b_z c_x d_y$$

$$a_x a_y b_z + a_x b_y d_z + b_x d_y d_z = a_y a_z b_x + a_z b_y d_x + b_z d_x d_y$$

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$$a_x d_x - b_x c_x = 1$$

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- The system is **independent** if it has no equivalent subsystem.
- **Big open question**
Is the size of an independent system of equations over n unknowns bounded?
- Open already for three unknowns (trivial for two).
- Unbounded in free **groups**.
- Lower bound $\Omega(n^4)$ (explicit system by Karhumäki and Plandowski 1996, Karhumäki and Saarela 2011).

Bounds on the size of independent systems for three unknowns

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Let E_1, \dots, E_m , $m \geq 2$, be an independent system of equations in three unknowns having a nonperiodic solution.

- Aleksi Saarela, Systems of word equations, polynomials and linear algebra: A new approach, European J. Combin. 2015

$$m \leq (|E_1|_x + |E_1|_y)^2 + 1 \text{ for some pair } x, y \text{ of unknowns.}$$

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- ŠH, Jan Žemlička, Algebraic properties of word equations, Journal of Algebra 2015

$$m \leq 2(|E_1|_x + |E_1|_y) + 1 \text{ for any pair } x, y \text{ of unknowns.}$$

Representation by polynomials

Let the alphabet A be a subset of \mathbb{N}_0 , and let unknowns be $\Theta = \{x, y, z\}$.

$$P : A^* \rightarrow \mathbb{N}_0[\alpha]$$

$$a_0 a_1 a_2 \cdots a_n \mapsto a_0 + a_1 \alpha + a_2 \alpha^2 + \cdots + a_n \alpha^n$$

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For a morphism $\varphi : \Theta^* \rightarrow A^*$, let

$$\mathcal{P}(\varphi) = (P(\varphi(x)), P(\varphi(y)), P(\varphi(z))) \in \mathbb{Q}(\alpha)^3$$

$$S : E \times \{x, y, z\} \rightarrow \mathbb{Z}[X, Y, Z]$$

$$E : (xyyz, zyyx)$$

$$S_{E,x} = 1 - ZY^2$$

$$S_{E,y} = X + XY - Z - ZY$$

$$S_{E,z} = XY^2 - 1$$

$$\mathcal{S}_E = (S_{E,x}, S_{E,y}, S_{E,z}) \in \mathbb{Q}(X, Y, Z)^3$$

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Length type $L = (L_x, L_y, L_z) \in \mathbb{N}_0^3$.

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Length type $L = (L_x, L_y, L_z) \in \mathbb{N}_0^3$. Define $\mathcal{S}_E(L)$ by morphism

$$\Omega_L : \mathbb{Z}[X, Y, Z] \rightarrow \mathbb{Q}(\alpha)$$

$$X \mapsto \alpha^{L_x} \quad Y \mapsto \alpha^{L_y} \quad Z \mapsto \alpha^{L_z}$$

φ with $L(\varphi) = \{|\varphi(x)|, |\varphi(y)|, |\varphi(z)|\}$ is a solution of E

if and only if

$$\mathcal{S}_E(L(\varphi)) \cdot \mathcal{P}(\varphi) = 0.$$

$$E_1 : (xyzz, zzxy)$$

$$E_2 : (xyxzz, zzxxy)$$

Representation by polynomials: Example (by Dirk Nowotka)

$$E_1 : (xyzz, zzxy)$$

$$E_2 : (xyxzz, zzxxy)$$

φ_i : a solution of the system **without** equation E_i

$$\varphi_{\emptyset} : \quad x \mapsto a \quad y \mapsto bab \quad z \mapsto ab$$

$$\varphi_1 : \quad x \mapsto a \quad y \mapsto babaabab \quad z \mapsto ab$$

$$\varphi_2 : \quad x \mapsto a \quad y \mapsto babab \quad z \mapsto ab$$

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If a common non-periodic solution φ , then $\mathcal{S}_{E_1}(L(\varphi))$ and $\mathcal{S}_{E_2}(L(\varphi))$ are linearly dependent over $\mathbb{Q}(\alpha)$.

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$$\varphi_\emptyset : \quad x \mapsto a \quad y \mapsto bab \quad z \mapsto ab$$

$$\mathcal{P}(\varphi) = (a, b + a\alpha + b\alpha^2, a + b\alpha)$$

$$\mathcal{P}(\varphi) = (1, \alpha, 1) \quad \mathcal{P}(\varphi) = (0, 1 + \alpha^2, \alpha)$$

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$$L(\varphi) = (1, 3, 2)$$

$$\mathcal{P}(\varphi) = (1, \alpha, 1) \quad \mathcal{P}(\varphi) = (0, 1 + \alpha^2, \alpha)$$

$$\mathcal{S}_{E_1}(L(\varphi)) = (1 - \alpha^4, \alpha - \alpha^5, \alpha^4 + \alpha^6 - 1 - \alpha^2)$$

$$\mathcal{S}_{E_2}(L(\varphi)) = (1 - \alpha^5, \alpha - \alpha^6, \alpha^5 + \alpha^7 - 1 - \alpha^2)$$

$$\mathcal{S}_{E_2}(L(\varphi)) = \frac{1 - \alpha^5}{1 - \alpha^4} \cdot \mathcal{S}_{E_1}(L(\varphi))$$

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GCD of all three 2×2 minors is

$$\tau(E_1, E_2) = X(1 + Z)(XY - Z^2).$$

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For what L , the polynomial τ vanishes under the mapping Ω_L ?

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$$L_x + L_y = 2L_z$$

$$|\varphi(x)| + |\varphi(y)| = 2|\varphi(z)|$$

Prize problem

I will pay **200 €** to the first person who gives the answer (with a proof) to the following question:

Is there a positive integer $n \geq 2$ and words u_1, u_2, \dots, u_n such that both equalities

$$\begin{cases} (u_1 u_2 \cdots u_n)^2 = u_1^2 u_2^2 \cdots u_n^2, \\ (u_1 u_2 \cdots u_n)^3 = u_1^3 u_2^3 \cdots u_n^3, \end{cases}$$

hold and the words u_i , $i = 1, \dots, n$, do not pairwise commute (that is, $u_i u_j \neq u_j u_i$ for at least one pair of indices $i, j \in \{1, 2, \dots, n\}$)?

Linear spaces of length types

Length types are not individuals; they form cones of dimension equal to the rank of the solution.

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Length types of solutions

$$\mathcal{L} = \{L(\vartheta_{k,\ell} \circ \varphi_{\emptyset}) \mid k, \ell \in \mathbb{N}\}$$

form a lattice in

$$\langle (1, 1, 1), (0, 2, 1) \rangle_{\mathbb{Q}}.$$

Vectors $\mathcal{S}_{E_1}(L)$ and $\mathcal{S}_{E_1}(L)$ are linearly dependent (i.e. the GCD of minors vanishes) for all $L \in \mathcal{L}$.

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Lemma

Let $\lambda \in \mathbb{Z}^n \setminus \{0\}$ have coprime coefficients and let $N \subseteq \mathcal{N}(\lambda)$ be of rank $n - 1$. Then $p \in \mathbb{Z}[\mathbf{X}]$ satisfies $\Omega_L(p) = 0$ for all $L \in N$ if and only if $(\mathbf{X}^{\lambda_\oplus} - \mathbf{X}^{\lambda_\ominus}) \mid p$.

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$$\lambda = (1, 1, -2) \quad \lambda_\oplus = (1, 1, 0) \quad \lambda_\ominus = (0, 0, 2) \quad (XY - Z^2)$$

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special form of \mathcal{S}_E



special form of τ



bound on the number of minimal monomials in τ



bound on the number of monomials $(\mathbf{x}^{\lambda_{\oplus}} - \mathbf{x}^{\lambda_{\ominus}})$ dividing τ

- Exploit better the special form of τ to improve the bound.
- Generalize the approach for all ranks, not just $n - 1$.