Towards computer-assisted semantic markup of mathematical documents

Year 1 progression talk

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June 13, 2023

## Introduction

- The language of mathematics (LoM) is the language used when producing mathematical documents
- Many attempts at computerising it
- Full of ambiguities!
- We usually typeset $P \times Q$ using $\$ P$ times $Q \$$
- Is this a matrix product or a Cartesian product?
- We can use STEX - a ATEX macro package for semantic markup - instead

- I hope to develop ways of automatically adding STEX
- This can improve learning materials as well


## The language of mathematics (LoM)

- The language used in mathematical documents
- In depth analysis by Ganesalingam and Iancu
- I will focus on the symbolic part - it has a clearer structure, and precise definitions so I'm hoping it will be easier to work with


## Attempts at computerising the language of mathematics

Representations of LoM

- Controlled natural languages (CNLs) - MathNat, ForTheL
- Discourse representation structures (DRS) - used in Ganesalingam's work
Software making use of these representations
- Proof assistants - Isabelle, Mizar, Coq, Lean, ...
- Proof checkers - Naproche, Naproche-SAD, ...


## Machine learning (ML)

- Could be used for processing text, lots of recent progress in natural language processing
- There have been attempts at using ML for working with LoM
- Müller and Kaliszyk tried to disambiguate symbolic formulas and using $\mathrm{S}^{\top} \mathrm{T}_{\mathrm{E}}$ to represent output
- Pagel and Schubotz tried to infer the meaning of symbols in a formula using surrounding text
- Hutterer tried generating suggestions for what $S^{\top} E X$ macros to use when typesetting formulas
- I lack experience with ML and would need to spend time learning it
- ML is probability-based so unsuitable on its own (even educated guesses aren't enough for mathematical truths)
- Could be useful for generating suggestions for the user


## Semantic markup

- Encoding the meaning of objects into the objects themselves
- Recall $P \times Q$ - \$P \times $Q \$$
- Meaning encoded within the formula - \$ $\backslash \operatorname{cart}\{P, Q\} \$$

In LATEX we can do this using STEX

## STEX — about

- A package for defining semantic macros in $\operatorname{LT}_{E} \mathrm{E} X$
- First released in 2008, but it was hard to use
- Difficult setup
- Complex internal dependency structure
- Major rework in 2022 to fix these issues
- Includes the Semantic Multilingual Glossary of Mathematics (SMGloM) - a library of macros for many areas of mathematics


## $S^{\top} E X$ - usage

Defining new macros

- We can write $a \times b$ as $\$$ a times $b \$$, but this can be ambiguous!
- Instead we can define an STEX macro \symdef\{realmult\}[args=2]\{\#1 \times \#2\}
- Now we can write $a \times b$ as \$\realmult $\{a\}\{b\} \$$

However...

## $S^{\top} E X$ - usage

This is not practical!

- What if we wish to write $a \times b \times c \times d$ ? Instead of using many \realmult calls, we use flexary arguments
- Redefine the macro to
\symdef \{realmult\}[args=a]\{\argsep\{\#1\}\{\times\}\}
- Now we can write $a \times b \times c \times d$ as $\$ \backslash r e a l m u l t\{a, b, c, d\} \$$

Other features

- Defining new notations -
\notation\{realmult\}[dot]\{\argsep\{\#1\}\{\cdot\}\} Using \$\realmult [dot] $\{\mathrm{a}, \mathrm{b}, \mathrm{c}, \mathrm{d}\} \$$ produces $a \cdot b \cdot c \cdot d$


## The $\lambda$-calculus - introduction

We use it as a testing ground for implementations

- Developed by Alonzo Church in the 1930s
- Turing-complete model of computation
- Relatively simple - only 3 "building blocks"
- This will not be a formal introduction


## The $\lambda$-calculus - basic definition

- We use $x, y, z$ to range over $\mathcal{V}::=v \mid x^{\prime}$
- $\mathcal{V}$ is countably infinite
- We use $A, B, C$ to range over $\Lambda::=x|A B| \lambda x . A$
- Each term is either a variable, application, or abstraction
- $(\lambda x .(x y)), z$, and $((x y) z)$ are all terms
- We use brackets to disambiguate the structure, but that can be difficult to read


## The $\lambda$-calculus - notational conventions

Notational conventions

- Parentheses around application can be dropped - $(A B)=A B$
- Application is left-associative - $((A B) C)=A B C$
- The scope of an abstraction extends as far to the right as possible
- Multiple "nested" abstractions can be written with just one $\lambda$ and dot - $(\lambda x .(\lambda y .(\lambda z .(A))))=(\lambda x y z .(A))$
The terms $x y z$ and $\lambda x$.xy as trees



## The $\lambda$-calculus - operations

We define two operations here (there are more)

- Substitution - $A[x:=B]$ replaces all free occurrences of $x$ in $A$ by $B$
- Example: $(\lambda x .(x y))[y:=z]$ gives $(\lambda x .(x z))$
- $\beta$-reduction - represents "computation"
- $(\lambda x .(A)) B \rightarrow A[x:=B]$
- Example: $(\lambda x .(x y)) z \rightarrow(x y)[x:=z]$ which gives ( $z y$ )


## An sTEX module for the $\lambda$-calculus

Motivation

- STEX macros for the $\lambda$-calculus did not exist before
- Potentially useful for the Foundations 1 course (a course exclusively about the $\lambda$-calculus)
- Learning experience for using $\mathrm{S}^{\top} \mathrm{EX}$

An STEX module for the $\lambda$-calculus - implementation
Implementation

- I implemented 3 macros: \app, \abs, \var
- Needed support for notational conventions
- I added 9 notations - some of them are used in tandem

$$
\begin{array}{l|l}
\text { \app [nb] }\{\mathrm{A}\}\{\mathrm{B}\} & A B \\
\text { \abs }[\mathrm{nb}]\{\mathrm{x}\}\{\mathrm{A}\} & \\
\text { \abs [nob] }\{\mathrm{x}\}\{\mathrm{A}\} & \lambda x \cdot A \\
\text { \abs [nib] }\{\mathrm{x}\}\{\mathrm{A}\} & \lambda x \cdot(A) \\
\text { \abs [nb-nodot] }\{\mathrm{x}\}\{\mathrm{A}\} & (\lambda x \cdot A) \\
\text { \abs [nib-nodot] }\{\mathrm{x}\}\{\mathrm{A}\} & (\lambda x A \\
\text { labs [nested] } \mathrm{x}\}\{\mathrm{A}\} & x . A \\
\text { \abs [nested-nodot] }\{\mathrm{x}\}\{\mathrm{A}\} & x A \\
\text { \abs [full] } \mathrm{x}\}\{\mathrm{A}\} & (\lambda x \cdot(A))
\end{array}
$$

\abs [nb-nodot] $\{\mathrm{x}\}\{\backslash$ abs [nested-nodot] $\{\mathrm{y}\}\{$ \abs [nested] $\{\mathrm{z}\}\{\mathrm{A}\}\}\}$ to typeset $\lambda x y z . A$

## An STEX module for the $\lambda$-calculus - redesign

There were issues

- It was impractical - many notations, needed 3 different ones for typesetting $\lambda x y z . A$
- \abs[nested-nodot] $\{\mathrm{x}\}\{\mathrm{y}\}$ and $\backslash \operatorname{app}[n b]\{\mathrm{x}\}\{\mathrm{y}\}$ look identical when typeset - could be ambiguous
I started redesigning the macros for better usability
- \abs now uses flexary arguments - \abs\{x, y, z\}\{A\}
- Unresolved issues with balanced brackets when using the "full notation" e.g., $(\lambda x .(\lambda y .(\lambda z .(A))))$
- Unfinished, but plan to finalize the changes soon


## A motivating example

Suppose you're writing a document on the $\lambda$-calculus

- Automatic removing of brackets, showing steps in substitutions and $\beta$-reductions
- Having it done automatically with macros inside $A T_{E X}$ would be convenient
Example

```
\begin{tabular}{l|l} 
\removeBrackets \(\{(\lambda x .(x y))\}\) & \(\lambda x . x y\)
\end{tabular}
\betaReduce \(\{(\lambda y z . z(y z))(\lambda x . x)\} \quad(\lambda y z . z)(y z)(\lambda x . x) \equiv_{\beta}\)
\(\lambda z . z((\lambda x \cdot x) z) \equiv_{\beta}\)
\(\lambda z . z(\lambda x . x) z \equiv_{\beta}\)
\(\lambda z . z z\)
```

I created something similar - lambda-calculus.lua

## lambda-calculus.lua - implementation

Implemented in Lua - easy interface with LATEX via macros Current functionality

- Converting $\mathrm{ST}^{\mathrm{E}} \mathrm{X}$ macros to trees
- Operating on trees is easier than on strings of macros
- Performing substitution and $\beta$-reduction on the trees
- Substitution returns just the end result of the computation
- $\beta$-reductions returns a list of all the steps
- Applying notational conventions
- This always tries to minimize the number of brackets that are needed and "compresses" nested abstractions
Note: it needs to be updated due to the $\lambda$-calculus macro redesign and $\mathrm{S}^{\top} \mathrm{E}_{\mathrm{E}}$ updates


## lambda-calculus.lua - discussion

- Requires the use of $S^{T} E X$ macros
- Showcases advantage of using semantic macros in documents
- Potential application in the Foundations 1 course
- Substitution and notational conventions aren't applied step by step (just the end result is returned)
- Upgrades - more $\beta$-reduction strategies, de Bruijn indices


## A grammar for parsing human-written $\lambda$-calculus

Motivation

- STEX is required to use lambda-calculus.lua
- Manual conversion of formulas to $\mathrm{S}^{\top} \mathrm{EX}$ macros takes time
- Automated parsing might be a solution

Found pyparsing, a parsing library for Python

- Comes with its own grammar syntax and parser
- I'm profficient in Python - easier to build stuff with pyparsing


## A grammar for parsing human-written $\lambda$-calculus

- Created a grammar that can parse human-written $\lambda$-terms
- Proof of concept, only works in a controlled environment
- Used it on small documents - created modified copies with "s ${ }^{T} \mathrm{EX}$-ified" formulas


## A grammar for parsing human-written $\lambda$-calculus

Discussion

- Creating the grammar took a long time
- It didn't accurately reflect the typesetting of a term, only its structure
- It only worked in a controlled environment (no subscripts in variable names, or handling of typesetting-only $I A T_{E} X$ code)
Not used in future work, also due to pyparsing drawbacks


## Switch from pyparsing

pyparsing proved inadequate for several reasons

- Not exhaustive - not good for ambiguous grammars
- No backtracking - parsing success depends on the order in which rules are applied
I searched for another Python library so I could reuse some code
- Focus on GLR-based libraries - exhaustive parsing
- Tried GLRParser
- No extensive documentation
- Terminals could not be capitalised (which is needed in mathematics)
- Found parglare


## parglare

- Supports LR and GLR parsing
- Comes with a grammar syntax with potentially useful extra functionality such as rule priorities and attaching meta-data
- Ignores whitespace in input sentences - like $T_{E} X$ 's math mode
- Extensive documentation and still updated

Made the switch when working on automatic grammar generation

## Automatic grammar generation

## Motivation

- Manual creation takes lots of time
- Lots of STEX macros already exist to represent output

Current idea

- Create a grammar from $\mathrm{S}^{\top} \mathrm{EX}$ macros
- Parse documents with it, produce $\mathrm{S}^{\mathrm{T}} \mathrm{EX}$ equivalents of formulas

Neither of these ideas are final and might change

## Automatic grammar generation - implementation

Supporting infrastructure

- Extracting formulas from .tex files
- Currently only \$s, will add support for other delimiters in the future
- Finding macro definitions in an $S^{T} E X$ module
- Currently focused on STEX definitions, support for user-defined macros is planned for the future
- Definition objects to store information
- Used for storing information about definitions extracted from macros
- Keeps track of the name, any notations, stores the ${ }^{4} T_{E X}$ source code and the grammar rules generated from it


## Automatic grammar generation - implementation

Rule generation

- Each argument placeholder is replaced by arg
- Each \argsep is turned into a separate rule of the form exlist $\rightarrow$ ex exlist|ex
Example:
\symdef\{abs\}[args=ai]\{\lambda \argsep\{\#1\}\{\} . \#2\} generates the following rules:

$$
\begin{aligned}
& \text { abs } \rightarrow \text { "\lambda" abs1list "." arg } \\
& \text { abs1list } \rightarrow \text { abs1 abs1list | abs1 }
\end{aligned}
$$

## Automatic grammar generation - implementation

Grammar generation and parsing

- Rules are converted to parglare syntax
- A rule is added for arg - each non-terminal can be on the right side (e.g., arg $\rightarrow$ abs)
- Rules with just one arg are converted into regular expressions instead
- Parsing small example sentences and finding all parses


## Automatic grammar generation - discussion

Drawbacks

- The grammars overgenerate - xyzw generates 8 parse trees, but only 1 is correct
- The arg rule is too broad - not everything can be an argument to everything else
- Regular expressions are temporary and unsuitable for arbitrary ATTEX
- Potentially we could use more information provided by macro definitions (types, precedence, associativity)
In the future we will likely need to work with a TEX parser to handle user-defined macros and other math-mode delimiters


## Future work

- Finishing the rework of the $S^{T} E X$ module for the $\lambda$-calculus
- Updating lambda-calculus.lua to work with the new $S^{\top} \mathrm{T}^{\mathrm{E}} \mathrm{X}$ module
- Creating a demonstration for interested users
- Figuring out improvements for all the mentioned drawbacks


## Questions

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