



SCHOOL OF MATHEMATICAL AND COMPUTER SCIENCES

Department of Mathematics

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F17LP

Logic and Proof

Semester 1 – 2019/20

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Duration: 2 Hours

Attempt all questions

A University approved calculator may be used  
for basic computations, but  
appropriate working must be shown to obtain full credit.

(Throughout this exam paper,  
*wff* is an abbreviation for *well formed formula(e)*)

1. (a) Construct truth-tables for  $p \wedge q$ ,  $p \vee q$ ,  $p \rightarrow q$ ,  $p \leftrightarrow q$  and  $p$  **nand**  $q$  [5 marks].
- (b) Construct the parse-tree of  $\neg(p \leftrightarrow q) \wedge (\neg p \rightarrow \neg r)$  [5 marks].
- (c) Construct the truth-table of  $\neg(p \leftrightarrow q) \wedge (\neg p \rightarrow \neg r)$  [5 marks].
- (d) Construct a wff in disjunctive normal form that has the following truth-table [5 marks].

$p$	$q$	$r$	$A$
$T$	$T$	$T$	$T$
$T$	$T$	$F$	$F$
$T$	$F$	$T$	$F$
$T$	$F$	$F$	$F$
$F$	$T$	$T$	$F$
$F$	$T$	$F$	$T$
$F$	$F$	$T$	$F$
$F$	$F$	$F$	$F$

- (e) Prove that  $\neg(p \vee (q \vee r))$  is logically equivalent to  $\neg p \wedge (\neg q \wedge \neg r)$  using truth tables. Justify your answer [5 marks].

**Exam continues ...**

2. (a) Prove that **nand** is an adequate set of connectives on its own. You may assume that  $\{\neg, \wedge\}$  is an adequate set of connectives. [5 marks].
- (b) Use **truth trees** to determine whether the following is satisfiable

$$(p \rightarrow q) \wedge (p \rightarrow \neg q)$$

[5 marks].

- (c) Use **truth trees** to write the following in disjunctive normal form

$$p \rightarrow (q \rightarrow p)$$

[5 marks].

- (d) Use **truth trees** to determine whether the following is a valid argument

$$p \vee q, \neg p \vee r \vdash q \vee r$$

[5 marks].

- (e) Use **truth trees** to determine whether the following is a tautology

$$(p \vee (q \wedge \neg q)) \leftrightarrow p$$

[5 marks].

**Exam continues ...**

3. (a) Prove that  $a^2 = a$  in a Boolean algebra. To gain full credit, you must make explicit reference to the Boolean algebra axioms listed at the end of this exam paper [5 marks].
- (b) Write down the input/output table of a transistor [1 mark]. Prove that not-gates and or-gates can be constructed from transistors alone [4 marks].
- (c) Write down a Boolean expression for the following input/output behaviour [2 marks] and so construct a corresponding circuit using not-gates, and-gates and or-gates [3 marks].

$x$	$y$	$z$	$u$
1	1	1	0
1	1	0	1
1	0	1	0
1	0	0	0
0	1	1	1
0	1	0	0
0	0	1	0
0	0	0	1

**Exam continues ...**

4. (a) Let  $R(x)$  be the predicate ' $x$  is a rabbit' and let  $L(x)$  be the predicate ' $x$  likes carrots'. Write down the first-order sentences 'all rabbits like carrots' and 'some rabbits like carrots' [5 marks].
- (b) Use **truth-trees** to prove that

$$(\exists x)[D(x) \rightarrow (\forall y)D(y)]$$

is universally valid [10 marks].

### Boolean algebra axioms

(B1)  $(x + y) + z = x + (y + z)$ .

(B2)  $x + y = y + x$ .

(B3)  $x + 0 = x$ .

(B4)  $(x \cdot y) \cdot z = x \cdot (y \cdot z)$ .

(B5)  $x \cdot y = y \cdot x$ .

(B6)  $x \cdot 1 = x$ .

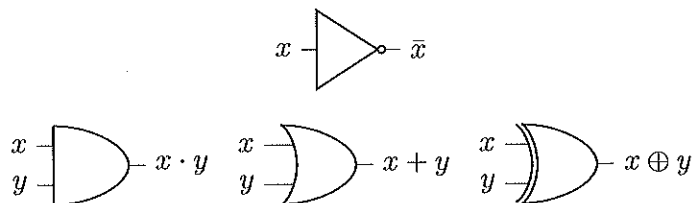
(B7)  $x \cdot (y + z) = x \cdot y + x \cdot z$ .

(B8)  $x + (y \cdot z) = (x + y) \cdot (x + z)$ .

(B9)  $x + \bar{x} = 1$ .

(B10)  $x \cdot \bar{x} = 0$ .

### Circuit symbols



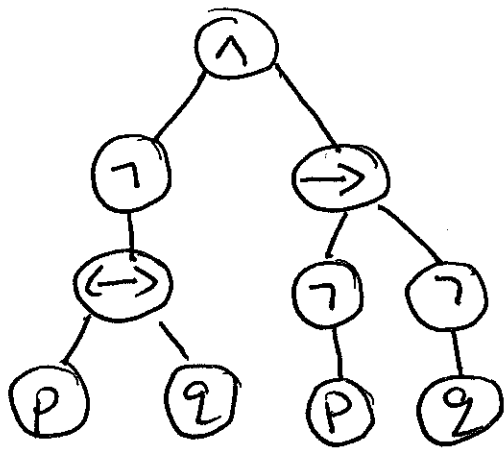
Exam ends

Logic and proof 2019 exam

1(a).

	[1]	[1]	[1]	[1]	[1]
P	Q	$P \wedge Q$	$P \vee Q$	$P \rightarrow Q$	$P \text{ NAND } Q$
T	T	T	T	T	F
T	F	F	T	F	T
F	T	F	T	T	T
F	F	F	F	T	T

(b)



[5]

(c)

P	Q	r	$\neg(P \leftrightarrow Q) \wedge (\neg P \rightarrow \neg r)$
T	F	T	T
T	F	F	T
F	T	F	T

[For brevity, I have just indicated where the outputs are T].

$$(d) (p \wedge q \wedge r) \vee (\neg p \wedge \neg q \wedge \neg r) \equiv A. [5]$$

(e)  $\neg(p \vee q \vee r)$  is only T when  $p, q, r$  are all F.  $\neg p \wedge \neg q \wedge \neg r$  is only T when  $p, q, r$  are all F. Since the truth tables are the same, we can deduce that

$$\neg(p \vee q \vee r) \equiv \neg p \wedge \neg q \wedge \neg r.$$


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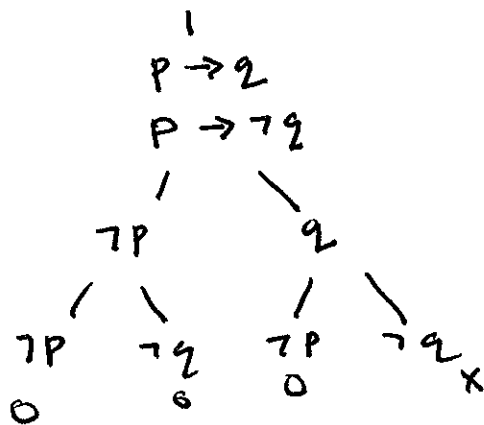
$$2(a) \quad p \text{ nand } q \stackrel{\text{def}}{=} \neg(p \wedge q)$$

$$\bullet \quad p \text{ nand } p = \neg(p \wedge p) \equiv \neg p. \quad [2]$$

$$\begin{aligned} \bullet \quad p \wedge q &\equiv \neg(\neg(p \wedge q)) \\ &\equiv \neg(p \text{ nand } q) \\ &\equiv (p \text{ nand } q) \text{ nand } (p \text{ nand } q). \quad [3] \end{aligned}$$

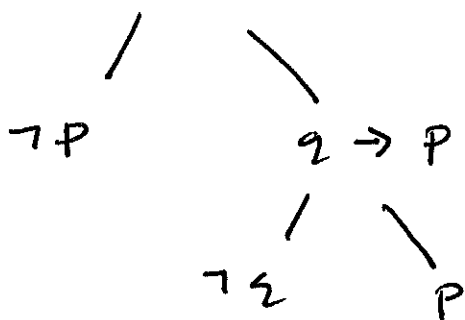

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$$(b) \quad (p \rightarrow q) \wedge (p \rightarrow \neg q) \checkmark$$



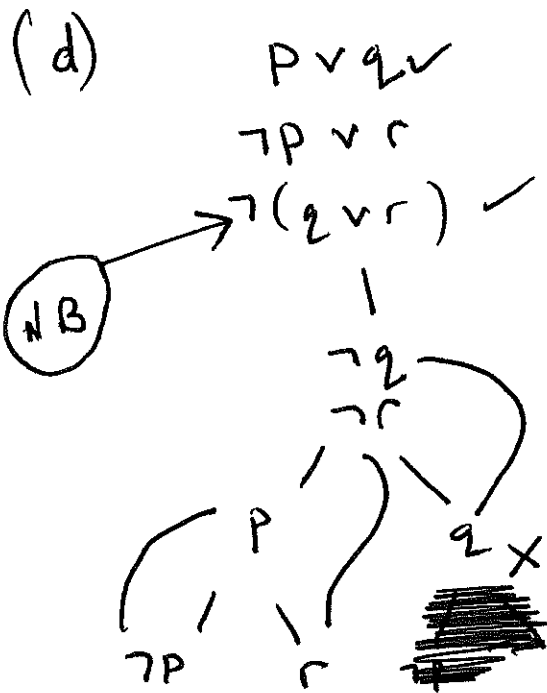
satisfiable since there is an open branch. [5]

$$(c) \quad p \rightarrow (q \rightarrow p)$$

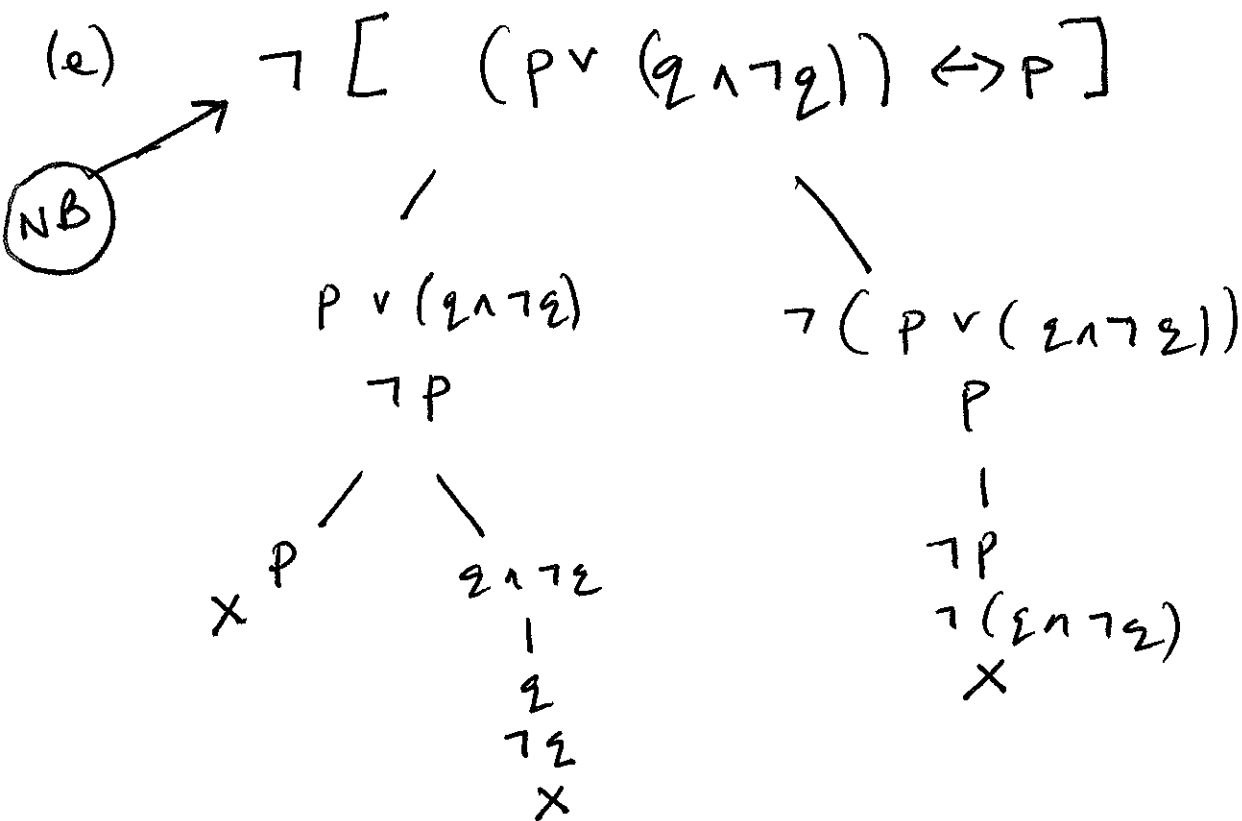


$$\therefore p \rightarrow (q \rightarrow p) \equiv (\neg p) \vee (\neg q) \vee (p). \quad [5]$$





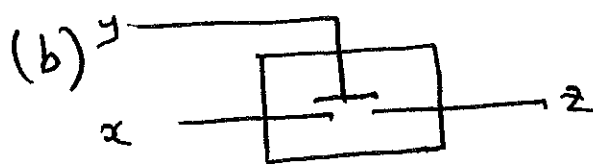
Tree closes and  $\therefore$  argument is valid.



Tree closes and  $\therefore$  the original iff is a tautology.

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$$\begin{aligned}
 3(a) \quad a &= a \cdot 1 \quad \text{by (B6)} \\
 &= a \cdot (a + \bar{a}) \quad \text{by (B9)} \\
 &= a \cdot a + a \bar{a} \quad \text{by (B7)} \\
 &= a^2 + 0 \quad \text{by (B10)} \\
 &= a^2 \quad \text{by (B3)} \quad [5]
 \end{aligned}$$



(schematic diagram of a transistor)

x	y	z
1	1	0
1	0	1
0	1	0
0	0	0

[1]

$$z = x \cdot \bar{y} = x \square y$$

my notation

•  $1 \square y = 1 \cdot \bar{y} = \bar{y}$ . Thus we can construct a not-gate. [2]

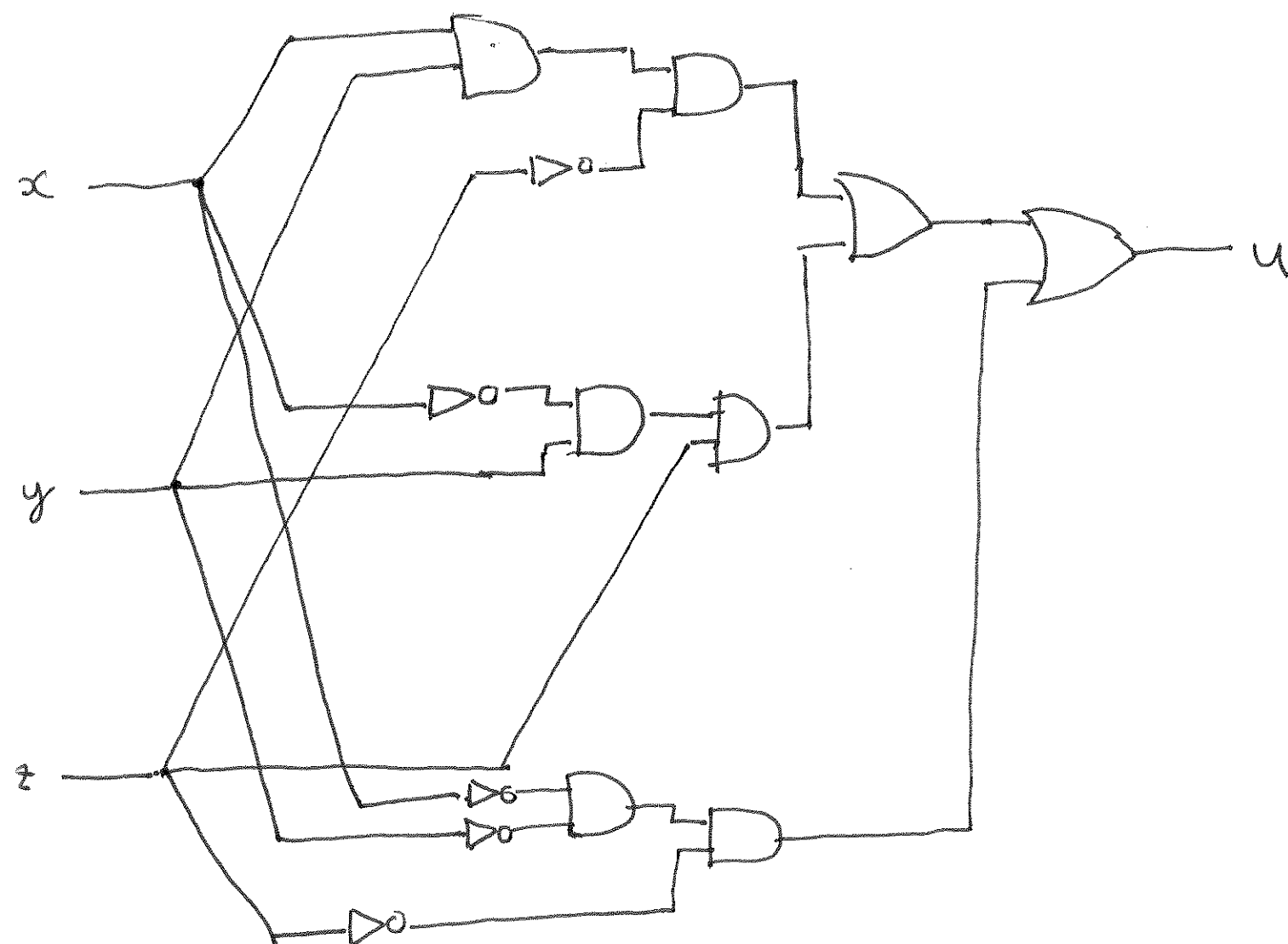
$$\begin{aligned}
 \cdot x + y &= \overline{\overline{(x + y)}} = \overline{(\bar{x} \cdot \bar{y})} = \overline{(\bar{x} \square y)} \\
 &= \overline{((1 \square x) \square y)} = 1 \square ((1 \square x) \square y).
 \end{aligned}$$

Thus we can construct an or-gate. [2].

3(c)

$$u = (x \cdot y \cdot \bar{z}) + (\bar{x} \cdot y \cdot z) + (\bar{x} \cdot \bar{y} \cdot \bar{z}) \quad [2]$$

[3]



4(a)

'all rabbits like carrots' =  $(\forall x)(R(x) \rightarrow L(x))$  [2½]'some rabbits like carrots' =  $(\exists x)(R(x) \wedge L(x))$  [2½]
 $\neg (\exists x) [D(x) \rightarrow (\forall y) D(y)] \quad \checkmark$ 

NB

 $(\forall x) \neg [D(x) \rightarrow (\forall y) D(y)] (*)$ 
 $\neg [D(a) \rightarrow (\forall y) D(y)] \quad \checkmark$ 
 $D(a)$ 
 $\neg (\forall y) D(y) \quad \checkmark$ 
 $(\exists y) \neg D(y)$ 
 $\neg D(b) \leftarrow \text{new name}$ 
 $\neg [D(b) \rightarrow (\forall y) D(y)]$ 
 $D(b)$ 
 $\neg (\forall y) D(y)$ 

[10]

X

Truth ~~tree~~ <sup>tree</sup> closes so wff is universally valid.